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Practical Quantum Key Distribution

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- QKD Protocols
- Implementations of QKD
- Photonic qubit QKD
- Channels
- Example: The Waterloo QKD





Review articles:

- N. Gisin et al., Quantum Cryptography, Reviews of Modern Physics 74, 145 (2002).
- M. Dusek et al., *Quantum Cryptography,* Progress in Optics **49,** 381 (2006).
- V. Scarani et al., *A Framework for Practical Quantum Cryptography,* arXiv:0802.4155, to appear in RMP.

Book:

 G. Van Assche, Quantum Cryptography and Secret-Key Distillation, Cambridge University Press (2006).





- Alice sends single photons with 1out-of-4 polarizations
- Bob measures in either + or × basis and gets one of two results (0, 1) in either case.
- Basis choices are announced after the measurement via authenticated public classical channel (internet, broadcast, ...)
- Measurement results for agreeing bases are key bits

C. H. Bennett & G. Brassard, *Quantum Cryptography: Public-key distribution and coin tossing* in *Proceedings of IEEE International Conference on Computer Systems and Signal Processing, IEEE*, 175-179 (1984).

The role of security proofs



- Security proofs give a lower bound on the achievable secure key rate as a function of **measurable** parameters
- They tell us how much key has to be sacrificed in privacy amplification in order to eliminate Eve's partial knowledge
- Shor & Preskill, PRL 85, 441 (2000): through reduction to entanglement purification and quantum error correction the secret key length is lower bounded by a factor of

1 - 2h(QBER) $h(x) = -x \log x - (1 - x) \log(1 - x)$

w.r.t the number of sifted bits, with exponentially small knowledge of the eavesdropper.

- Therefore if QBER < 11%, the secret key length is finite.
- With imperfect error correction we need to use

 $1 - h(QBER) - h_{EC,leakage}$

Noisy Keys



- Raw keys are noisy, because of errors in
 - Channel
 - Equipment (dark counts)
 - Eavesdropper
 - (Classical) Error correction can eliminate errors
 - Simple example: Take two blocks of k bits, compare parity, if different, dicscard





Alice new 10001011 11011110 00110011 11001111 11010110 01010 Bob new 10001011 11011110 00110011 11001111 11010110 01010

Simplified Cascade Error Correction

- Optimized for computational efficiency (vs. information leakage)
- 4 passes
 - Use QBER as determined in previous chunk to choose block size
 - Split key into blocks (randomly chosen bit order, different for each pass)
 - Apply BINARY to correct one error in each block (for odd numbers of errors)
 - Calculate parity
 - On disagreeing parity divide block in half
 - Repeat until error found
 - If error is found in later pass, there must have been even number of errors in previous pass' block → go back and correct using BINARY
 - Keep track of every bit sent via the public channel



BER	Simplified Cascade	Full Cascade
0.01	0.089	0.085
0.025	0.197	0.1925
0.05	0.341	0.335
0.075	0.477	0.465
0.1	0.589	0.577
0.125	0.717	0.697
0.15	0.817	0.805

G. Brassard and L. Salvail, "Secret-Key Reconciliation by Public Discussion," Advances in Cryptology – EUROCRYPT '93, LNCS 765, 410 (1994).

Privacy Amplification

- All the bits revealed during error correction must be discarded
- Any information an eavesdropper could have according to the QBER can be made exponentially small by hashing
- Determine final key length estimate: $R = N (1 - h_2(QBER)) - #(bits leaked) - #(security bits)$
- Shor-Preskill: R = N 2H₂(QBER)
 Since #(bits leaked) > H₂(QBER)
 this is always secure
- Calculate $k = (m * (raw key) + n) \mod p$
 - *m, n* are random number generated from a shared seed
 - *p* is a shared big prime number
- Use the last *R* bits of *k* as the key





Privacy Amplification



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By source

Prepare and measureEntanglement based

By Modulation

- Discrete
- Continuous
- Distributed phase reference

By implementation

- One-way
- Plug & Play



Discrete = Qu**d**its

- Polarization
- Time-bin
- Spatial Mode

Continuous Variables

Quadratures of field modes

Distributed Phase Reference

- Differential Phase Shift
- Coherent One-Way





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Entanglement Based





- Source can be under eavesdroppers control
- Immune to sidechannels

Continuous Variables



- Alice sends coherent states with a random modulation in a given quadrature
- Bob chooses randomly to measure a quadrature using homodyne detection
- Alternative: Squeezed states

$$\mathbf{E}_{\mathbf{k}} = i \sqrt{\frac{\hbar\omega_{\mathbf{k}}}{2\epsilon_{0}}} \left[\hat{a}_{\mathbf{k}} \mathbf{u}_{\mathbf{k}} e^{-i\omega_{\mathbf{k}}t} - \hat{a}_{\mathbf{k}}^{\dagger} \mathbf{u}_{\mathbf{k}} e^{i\omega_{\mathbf{k}}t} \right]$$

$$|lpha
angle = e^{-|lpha|^2/2} \sum_{n} \frac{lpha^n}{\sqrt{n!}} |n
angle$$

Continuous Variables







Distributed Phase Reference



D. Stucki et al., APL **87**, 194108 (2005).



Photon Polarization

- Every mode has two orthogonal polarizations (directions of the electric field)
- Arbitrary polarization states are superpositions
- Classically, polarization is described on the Poincaré sphere





The Dual Rail Qubit







For stability one can multiplex the two rails onto one.







Sources

• Attenuated lasers: poissonian statistics

$$p(n,\mu)=\frac{\mu^n}{n!}e^{-\mu}$$

• In order to optimize the secure key rate μ has to be set to a value that scales with *t*, the transmission of the channel

$$\mu_{\rm opt} \approx t\eta \frac{1 - h(\text{QBER}) - h(2 \text{QBER})}{1 - h(2 \text{QBER})}$$

$$K \approx R \frac{1}{2} \mu_{\text{opt}} \left[1 - h(\text{QBER}) - h(2 \text{QBER}) \right]$$







Single Photon



$K \approx R \left[1 - 2h(\text{QBER})\right]$

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0.25

0.50

0.75

Single Photons





Second Order Degree of Coherence

Measured by Hanbury Brown – Twiss Interferometry

$$g^{(2)}(\tau) = \frac{\langle I(t)I(t+\tau)\rangle}{\langle I(t)\rangle\langle I(t+\tau)\rangle}$$

CW

Two-photon suppression limited by detector resolution and finite reexcitation probability

Pulsed

Two-photon suppression only limited other background (filter performance)



 Alice randomly chooses from a few (e.g. 3) different mean photon numbers

$$\mu_{\text{opt}} \approx \frac{1}{2} \left[1 - \frac{h(\text{QBER})}{1 - h(\text{QBER})} \right] \qquad K \approx R \frac{1}{2} \mu_{\text{opt}} \left[1 - 2h(\text{QBER}) \right]$$

Modulation



- Combine multiple lasers and pulse them individually
 - Beware of side channels!

- Modulate laser
 - Polarization
 - Phase (commercially up to 40 GHz)
 - Amplitude for decoy



Some Facts About Detectors

- Detection efficiency = Quantum efficiency * Amplification efficiency
- For red / very near infrared light about 70%, ~10 /s noise
- Most common: Single Photon Avalanche Diode (SPAD)
- For telecommunication wavelengths (1550 nm): InGaAs APDs have <15% efficiency, some 10000 /s noise counts
- Alternative detectors
 - Visible Light Photon Counter
 - Superconducting Transition Edge
 Detector







Performance comparison





V. Scarani et al., arXiv:0802.4155

CVContinuous VariablesWCPWeak Coherent PulsesCOWCoherent One-WayEBEntanglement BaseddecoyDecoy States1-phSingle Photon Source

Platform BB84, COW	Parameter μ mean intensity V visibility: P&M V visibility: EB t_B transmission in Bob's device η det. efficiency p_d dark counts ϵ (COW) bit error ζ (EB) coherent 4 photons 0 0 leak EC code	Set #1 (opt.) 0.99 0.96 1 0.1 10 ⁻⁵ 0.03 0 1.2	Set #2 (opt.) 0.99 0.99 1 0.2 10 ⁻⁶ 0.01 0 1
CV	v= v _A + 1 variance	(opt.)	(opt.)
	ε optical noise	0.005	0.001
	η det. efficiency	0.6	0.85
	v _{el} electronic noise	0.01	0
	β EC code	0.9	0.9

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Guide light in singlemode optical fibers

Pfennigbauer et al., JON 4,





GAP-Optique, U. Geneva

Broadcast photons from a satellite using telescopes

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Optical Fibers

- Fused silica core guides light
- Attenuation by Rayleigh scattering
- Minimum @1550 nm: 0.17 dB/km = 4%/km loss
- Installed fiber typically has 0.3 dB/km
 - Polarization
 - Birefringence needs to be compensated
 - Depolarization due to different group velocities ($\sim \sqrt{L}$)





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Free-Space Optical Links



- Send photons through air in "beam"
- Diffraction causes beam to spread (~L²)
- Turbulence causes beam wander
 - → Can be incorporated as additional diffraction
- Scattering causes exponential attenuation

$$A = \frac{L^2(\theta_T^2 + \theta_{\text{atm}}^2)}{D_R^2} 10^{\frac{A_{\text{atm}}}{10}}$$





G. Bianco: The Matera Laser Ranging Observatory System

The MLRO telescope





- Diffraction angle ~(wavelength/diameter)
- Need stable pointing
- For satellites: tracking

Atmosphere



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Satellites

36000km



- From 1000 km altitude the horizon is 3000 km away
- Atmospheric attenuation becomes negligible above 10km



LEO satellites move fast

- Can only be "seen" from a ground stations for a small fraction of the orbit
- Diffraction loss becomes very severe for geostationary satellites

Early Experimental QKD

- 1989 Bennett et al., J. Cryptolog. 5, 3 (1992)
 30cm faint laser pulses
- 1993 Muller et al., Europhys. Lett. 23, 383 (1993)
 Polarization in fiber
- 1994 Townsend, Electron. Lett. **30**, 809 (1994)
 10 km fiber, phase
- 1996 Muller et al., Appl. Phys. Lett. 70, 793 (1997)
 Plug & play system
- 1999 Jennewein et al., Phys. Rev. Lett. 84, 4729 (2000)
 Entanglement based QKD (360m)
 1999 Tittel et al. Phys. Rev. Lett. 84, 4737 (2000)
 Energy-time entanglement in fiber



The plug & play system (67km demo)





Figure 1. Picture of the p&p system.



Figure 2. Schematic of the p&p prototype.

Stucki, et al., NJP **4,** 41 (2002).

- Uses phase encoding
- Eliminates polarization correction by Faraday mirror
- Need to send "strong" pulse from Bob to Alice for coding

Increasing the distance



Differential Phase Shift Keying QKD



Takasue et al., NJP **7**, 232 (2005). Diamanti et al., quant-ph/0608110.

- Better use of clock period
- Achieved 1 GHz clock rate
- Using up-conversion single photon detectors
- @100 km 166 bits/s secure (?)

Polarization in Fiber



Peng et al., quantph/0607129 (2006)

With decoy states achieved 103 km



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With (Almost) Noise-Free Detectors

 Superconducting Transition Edge Sensors

- Virtually zero noise
- Poor timing → slow clock cycle
- With decoy states achieved unconditionally secure key over 107 km

Rosenberg et al., Appl. Phys. Lett. **88**, 021108 (2006). Rosenberg et al., quant-ph/0607186

With Single Photons

- Single photons: unconditional security without decoy states
- Waks et al.: InAs quantum dots
- Alleaume et al.: Color centers in diamond

High Data-Rate Free-Space QKD

• 690 kbit/s at 0.15 photons/pulse at Alice

Free-Space Long Distance

- Entanglement-based with source at Alice's
- 1m receiver telescope
- Typical loss -30 dB
- ~30 raw key bits/s

Entanglement based FS QKD

- Dedicated real time entanglement based QKD system
- 630 bits/s *final* key

Marcikic et al., Appl. Phys. Lett. **89,** 101122 (2006)

IQC

CEIT

ΡΙ

The IQC-Perimeter Institute QKD Experiment

http://maps.google.com/maps/ms?msa=0&msid= 103964276287441386699.00000113448f5481181e2

Entangled Photon Pairs

Polarization entangled photon pairs via special geometry

Entangled Photon Pair Source

Send / Receive Equipment

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Bob @ Perimeter Institute

Alignment Spots

QBER

Time of Day

Key Rate

Corrected Key

		Alice				
		Н	V	+	-	
Bob	Н	39,497	1,218,454	393,100	355,074	2,006,125
	V	1,300,749	112,793	682,595	854,848	2,950,985
	+	680,032	878,628	51,217	1,262,143	2,872,020
	-	548,695	955,146	1,374,648	63,261	2,977,750
		2,604,973	3,165,021	2,501,560	2,535,326	

- Raw key rate = 565 bits/sec
- Sifted key rate = 284 bits/sec
- Optimum final secret key rate = 124 bits/sec
- Actual final secret key rate = 85 bits/sec
- QBER = 4.92%
- Total key of 1,612,239 bits > 1.5MB generated
- Visibilities: H/V = 88.6%, +/- = 91.7%
- Residual error rate = 1.92 e-003 errors/bit